

ECE 09468/09568

Discrete Event Systems

Lecture 10: **Petri Nets (PNs) – Part VI**

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Timed Petri Nets

- **Deterministic Timed PN (DTPN):**
 - Concurrent *choice-free* PN that associate *fixed time delays* to transitions, places, or arcs.
- **Stochastic Petri Nets (SPN):**
 - Time delays are associated with *transition only*
 - Time delay for each transition is assumed to be random and *exponentially distributed*.

What is Exponential Distribution

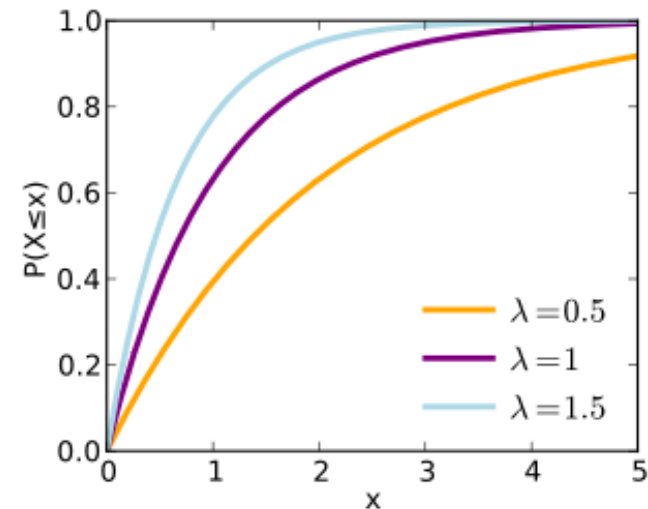
Given a random variable τ , if the probability that τ is greater than a given value x satisfies:

$$P(\tau > x) = e^{-\lambda x} \quad \text{for } x \geq 0$$

i.e. distribution function $F(x) = P(\tau \leq x) = 1 - e^{-\lambda x}$ Where $\lambda > 0$, called a rate.

τ is an **exponential random variable**.

If arrival intervals are such random variable, this arrival process is called **Poisson arrival process**.



Exponential Random Variables

If τ is an exponential random variable, then the mean (expected value) of τ is $E(\tau)=1/\lambda$ (Physically mean=average, $\lambda=1/\text{average}$ is thus called **rate**)

A transition in a PN can be associated with an exponentially distributed random time delay. In this case, λ is called **firing rate**.

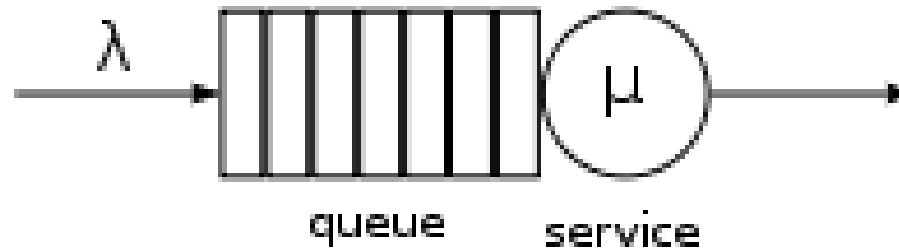
SPN Model

5-tuple : $SPN = (P, T, I, O, \Lambda, m_0)$

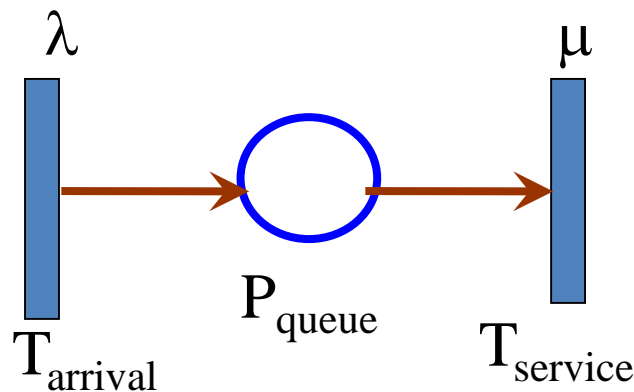
- P - A finite set non-empty of places
- T - A finite set non-empty of transitions
- I - A set of input arcs connecting P to T
- O - A set of output arcs connecting T to P
- $\Lambda: T \rightarrow R^+$ - Transition firing rates
- M_0 - Initial marking

M/M/1 Queue – PN representation

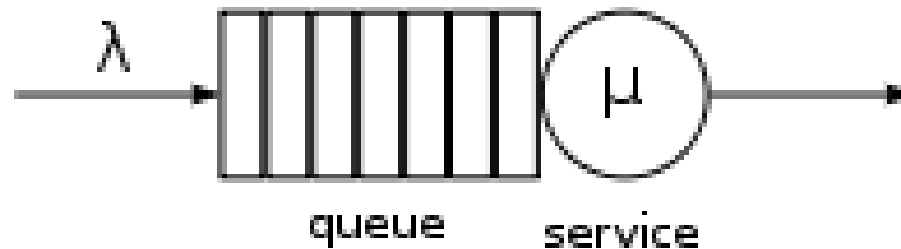
The M/M/1 is a queue with exponential arrivals and service times and a single server.



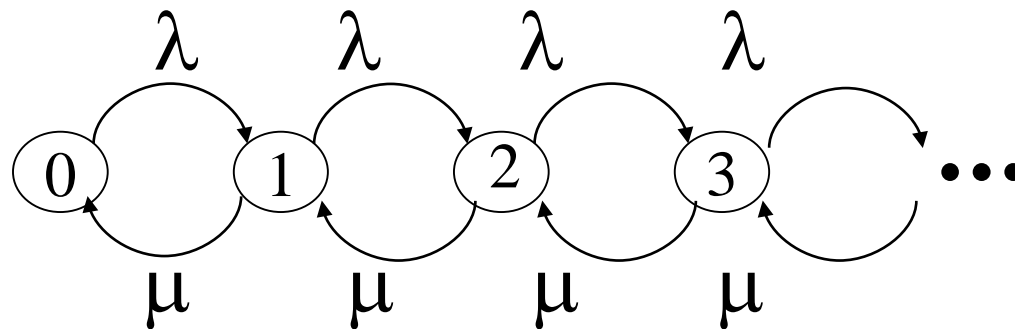
PN representation



Markov Chain Model of M/M/1



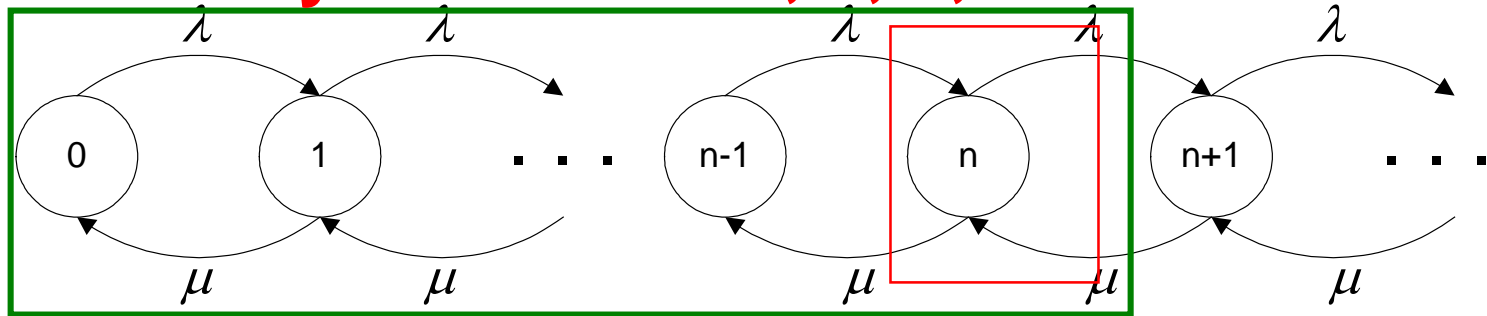
A **Markov chain** is a random process with the property that the next state depends only on the current state.



How to solve a Markov model for steady state probabilities, average wait time, and throughput?

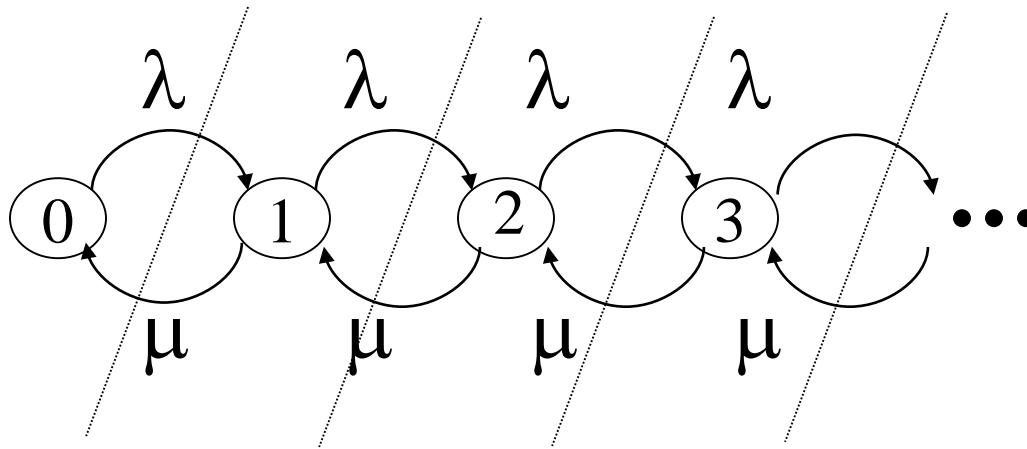
Markov Chain Model of M/M/1

- **How to find the π_i is the steady state probability for state $i=0, 1, 2, \dots$?**



- **Balance equation:** $(\lambda + \mu)\pi_n = \lambda\pi_{n-1} + \mu\pi_{n+1}$
- **Global balance equation:** $\lambda\pi_n = \mu\pi_{n+1}$
- **Normalization condition:** $\sum_{n=0}^{\infty} \pi_n = 1$

Solving Markov Model of M/M/1



List the balance equations (the probability flow from state n to state m equals the probability flow from state m to state n):

$$\mu\pi_1 = \lambda\pi_0, \mu\pi_2 = \lambda\pi_1, \mu\pi_3 = \lambda\pi_2, \dots$$

Solving Markov Model of M/M/1

Assume that π_0 is known and let $\rho = \lambda/\mu$, we find:

$$\mu\pi_1 = \lambda\pi_0 \Rightarrow \pi_1 = (\lambda/\mu)\pi_0 = \rho\pi_0$$

$$\mu\pi_2 = \lambda\pi_1 \Rightarrow \pi_2 = (\lambda/\mu)\pi_1 = \rho\pi_1 = \rho^2\pi_0$$

$$\mu\pi_3 = \lambda\pi_2 \Rightarrow \pi_3 = (\lambda/\mu)\pi_2 = \rho\pi_2 = \rho^3\pi_0$$

...

Since the sum of all probabilities is one, we have:

$$\pi_0 + \pi_1 + \pi_2 + \pi_3 + \dots = 1$$

$$\pi_0(1 + \rho + \rho^2 + \rho^3 + \dots) = 1$$

Solving Markov Model of M/M/1

$1 + \rho + \rho^2 + \rho^3 + \dots$ is a geometric series

Stability condition for M/M/1: $\rho < 1$ (service faster than arrival). Then we have:

$$\pi_0 = 1 - \rho$$

$$\pi_1 = \rho(1 - \rho)$$

$$\pi_2 = \rho^2(1 - \rho)$$

$$\pi_3 = \rho^3(1 - \rho)$$

...

Performance of M/M/1

If $\rho < 1$ (service faster than arrival), the throughput is

$$\begin{aligned}\gamma &= \mu\pi_1 + \mu\pi_2 + \mu\pi_3 + \dots \\ &= \mu(\pi_1 + \pi_2 + \pi_3 + \dots) = \mu(1 - \pi_0) = \mu\rho = \lambda\end{aligned}$$

Average number of customers in the system:

$$E(N) = \pi_1 + 2\pi_2 + 3\pi_3 + \dots = \rho / (1 - \rho)$$

Little's formula: $E(N) = \lambda E(T)$ where $E(T)$ is average delay

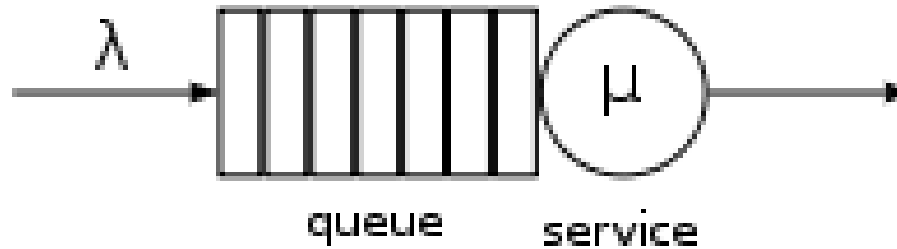
$$E(T) = E(N) / \lambda = [1 / \mu] / (1 - \rho) \text{ and}$$

Average wait: $E(W) = E(T) - 1 / \mu$

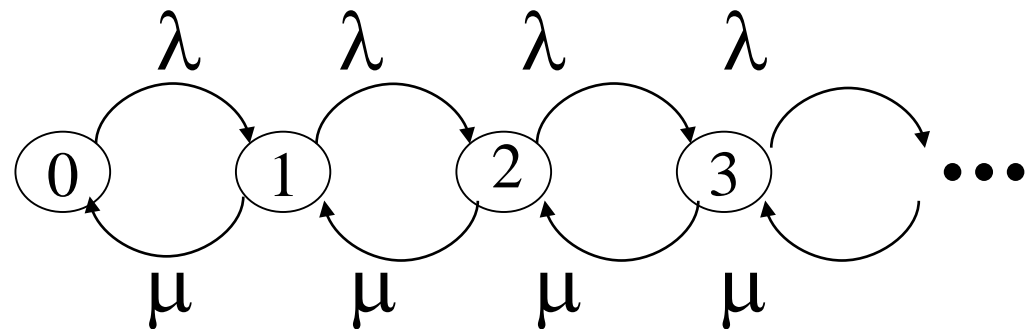
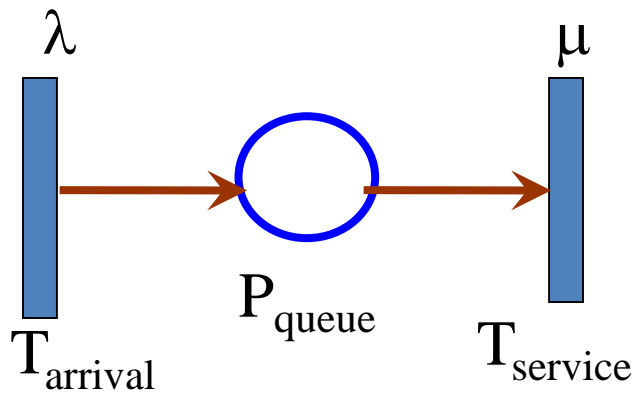
SPN (Stochastic Petri Net)

- SPNs are obtained by associating stochastic and timing information to Petri nets
 - Time delays are associated with *transition only*
 - Time delay for each transition is assumed to be random and *exponentially distributed*.

SPN model of M/M/1 Queue

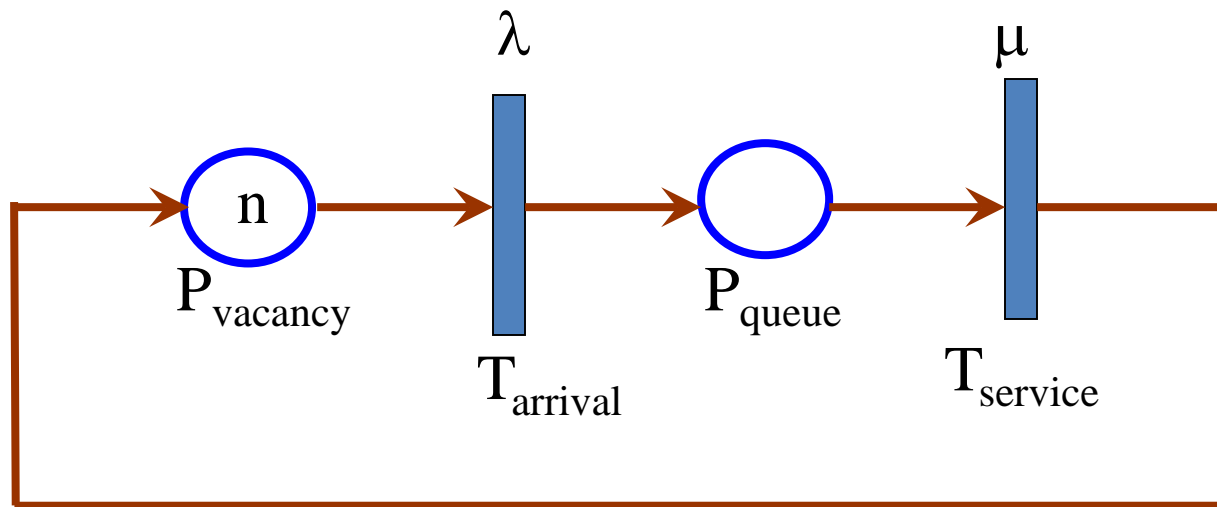


SPN representation



PN model of M/M/1/n Queue

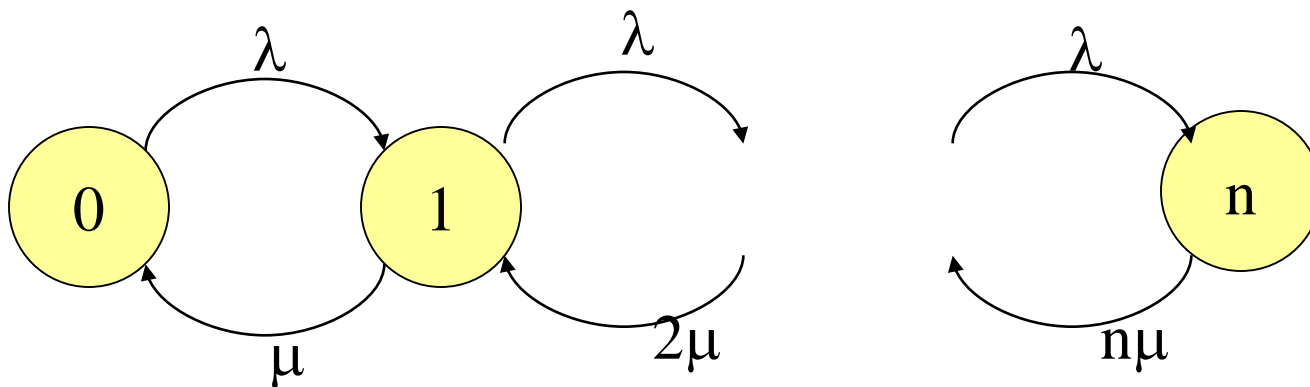
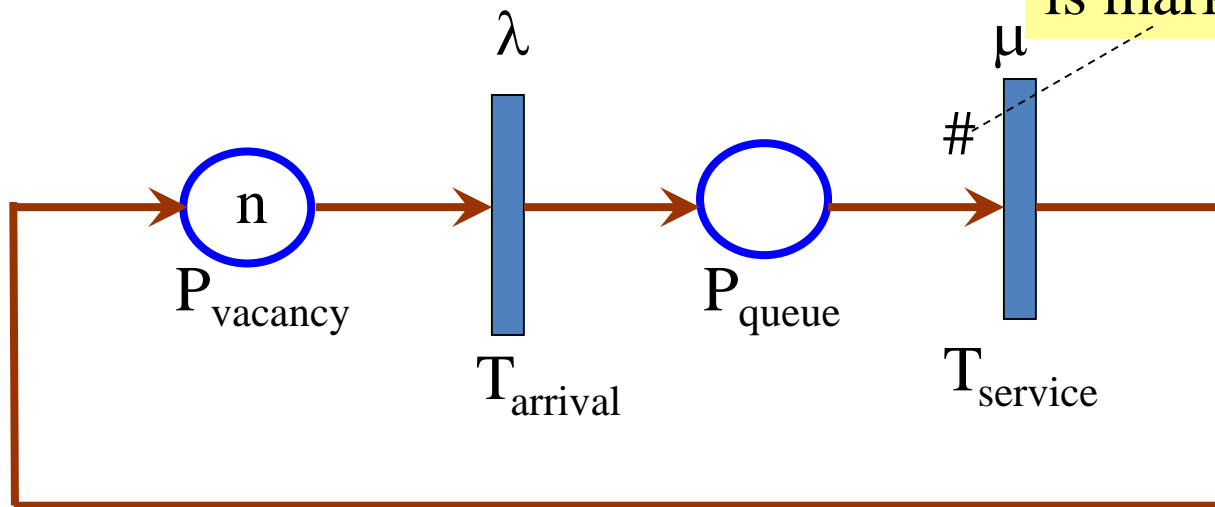
- M/M/1/n queue: an M/M/1 queuing system with a limited buffer space so that **at most n jobs** can be in the system at a time



What is the corresponding Markov model?

PN model of M/M/n/n Queue

- M/M/n/n queue: n servers and buffer size is n .



Solution Procedure for SPN

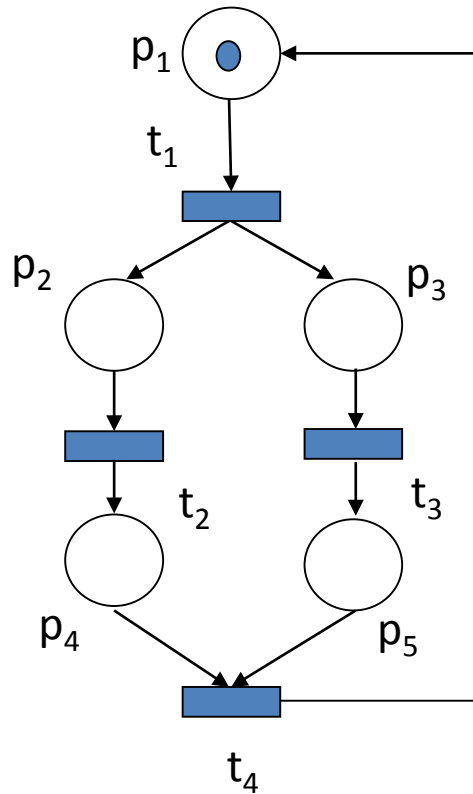
- Build an SPN including proper time delay variables characterized by firing rates.
- Find reachability graph labeling with firing rates.
- Build up a transition rate matrix A and solve $\pi A=0$ and $\sum \pi_i=1$ where

$$\pi = (\pi_0, \pi_1, \dots, \pi_{q-1})$$

Or list $q-1$ balance equations instead of $\pi A=0$

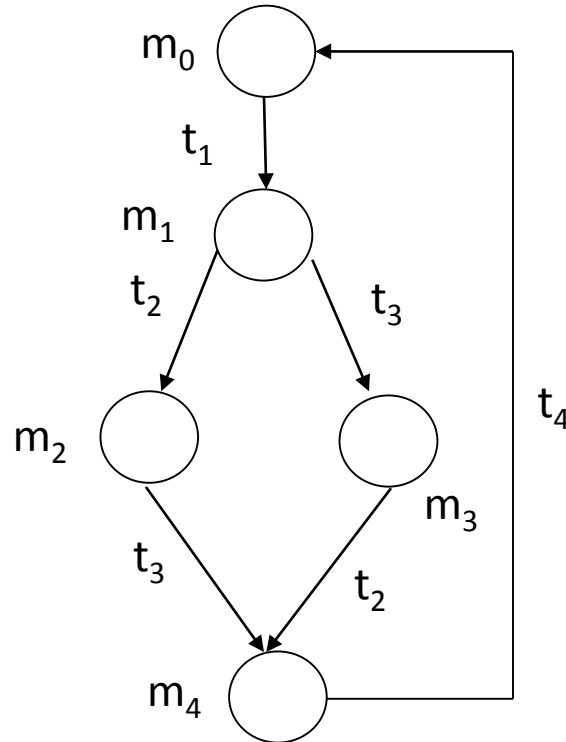
SPN example

Generating Reachability Graph:



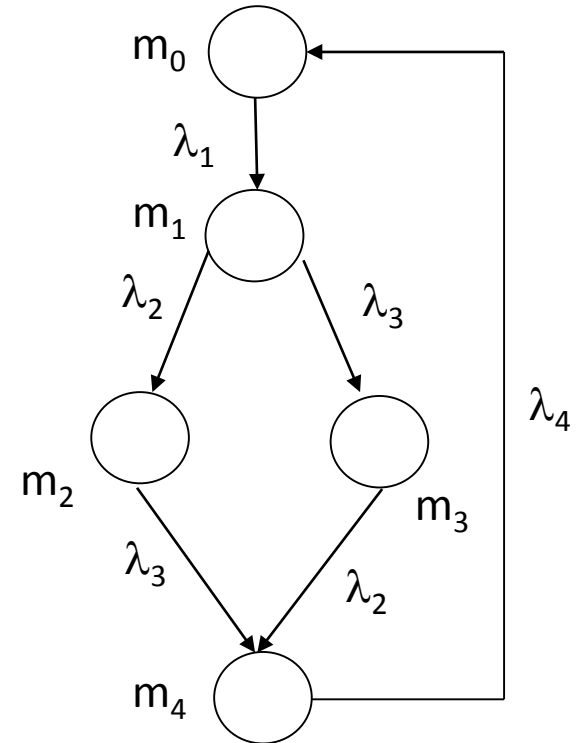
$$m_0 = (1\ 0\ 0\ 0\ 0)^T$$

$$m_3 = (0\ 1\ 0\ 0\ 1)^T$$



$$m_1 = (0\ 1\ 1\ 0\ 0)^T$$

$$m_4 = (0\ 0\ 0\ 1\ 1)^T$$



$$m_2 = (0\ 0\ 1\ 1\ 0)^T$$

SPN example

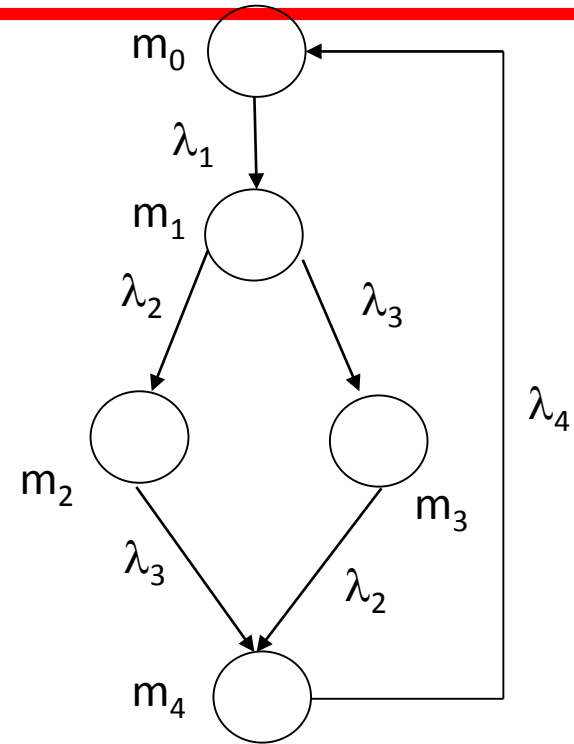
Transition rate matrix $A=?$

$$\begin{pmatrix} a_{00} & a_{01} & a_{02} & a_{03} & a_{04} \\ a_{10} & a_{11} & a_{12} & a_{13} & a_{14} \\ a_{20} & a_{21} & a_{22} & a_{23} & a_{24} \\ a_{30} & a_{31} & a_{32} & a_{33} & a_{34} \\ a_{40} & a_{41} & a_{42} & a_{43} & a_{44} \end{pmatrix}$$

Decide **non-diagonal** elements first.

Then find diagonal elements according to $a_{i0}+a_{i1}+a_{i2}+a_{i3}+a_{i4}=0$, $i=0, 1, 2$, and 3 .

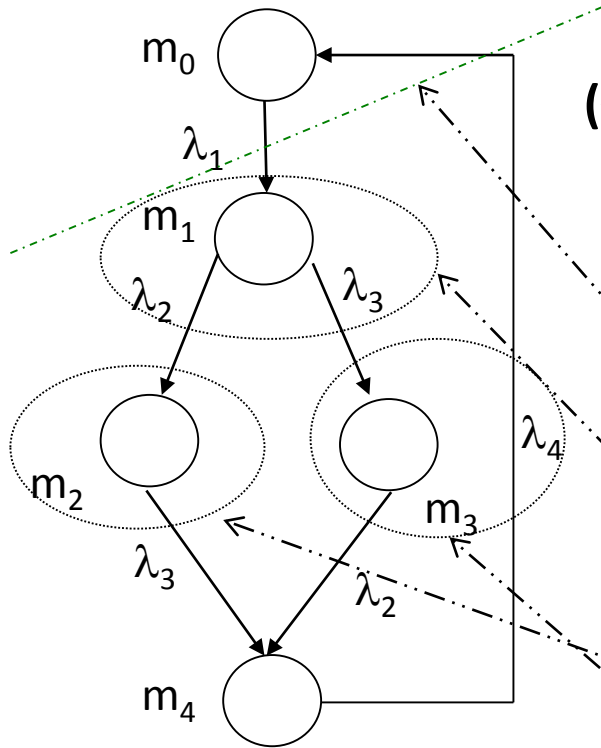
For $i \neq j$, a_{ij} =the total rate leaving from m_i to m_j .



$$\begin{pmatrix} -\lambda_1 & \lambda_1 & 0 & 0 & 0 \\ 0 & -(\lambda_2+\lambda_3) & \lambda_2 & \lambda_3 & 0 \\ 0 & 0 & -\lambda_3 & 0 & \lambda_3 \\ 0 & 0 & 0 & -\lambda_2 & \lambda_2 \\ \lambda_4 & 0 & 0 & 0 & -\lambda_4 \end{pmatrix}$$

SPN example

Balance equations



$(\pi_0, \pi_1, \pi_2, \pi_3, \pi_4)$

$$\begin{pmatrix} -\lambda_1 & \lambda_1 & 0 & 0 & 0 \\ 0 & -(\lambda_2 + \lambda_3) & \lambda_2 & \lambda_3 & 0 \\ 0 & 0 & -\lambda_3 & 0 & \lambda_3 \\ 0 & 0 & 0 & -\lambda_2 & \lambda_2 \\ \lambda_4 & 0 & 0 & 0 & -\lambda_4 \end{pmatrix}$$

$$-\lambda_1 \pi_0 + \lambda_4 \pi_4 = 0$$

$$\lambda_1 \pi_0 - (\lambda_2 + \lambda_3) \pi_1 = 0$$

$$\lambda_2 \pi_1 - \lambda_3 \pi_2 = 0$$

$$\lambda_3 \pi_1 - \lambda_2 \pi_3 = 0$$

SPN example

Produce Solution of Steady State Probabilities

$$-\lambda_1\pi_0 + \lambda_4\pi_4 = 0$$

$$\lambda_1\pi_0 - (\lambda_2 + \lambda_3)\pi_1 = 0$$

$$\lambda_2\pi_1 - \lambda_3\pi_2 = 0$$

$$\lambda_3\pi_1 - \lambda_2\pi_3 = 0$$

$$\pi_0 + \pi_1 + \pi_2 + \pi_3 + \pi_4 = 1$$

given $\lambda_1 = \lambda_2 = \lambda_3 = \lambda_4 = 1$,

$$\pi_0 = \pi_4 = 2/7$$

(The system is at either m_0 or m_4 with probability $2/7$)

$$\pi_1 = \pi_2 = \pi_3 = 1/7$$

The system is at either m_1 , m_2 , or m_3 with probability $1/7$)

- 1) Assume π_0 is known.
- 2) Represent others in terms of π_0 only.
- 3) Use the last equation to find π_0
- 4) Find π_1 to π_4 .

Performance Measures

1) Probability of a certain condition V being true.

Suppose $G = \{m: \text{The condition } V \text{ is true at } m \text{ in reachability graph}\}$. Then

$$\text{Prob}(V) = \sum_{m_i \in G} \pi_i$$

Examples:

A machine utilization, i.e., a machine is in use;

A machine failure rate, i.e., a machine is in failure;

A channel idle rate, i.e., a channel is idle;

...

Performance Measures

2) Average number of tokens in a place p.

Define condition V1: a token in p; V2: 2 tokens in p; ...; Vk: k tokens in p. Assume p has at most k tokens.

$$E(m(p)) =$$

$$\text{Prob}(V1) + 2\text{Prob}(V2) + \dots + k\text{Prob}(Vk)$$

$$= \sum_{i=1}^k i \times \text{Prob}(Vi)$$

Performance Measures

3) **Frequency of firing a transition**, i.e., the average number of times the transition fires in unit time. Let $G = \{m : t_j \text{ is enabled at } m \text{ in reachability graph}\}$.

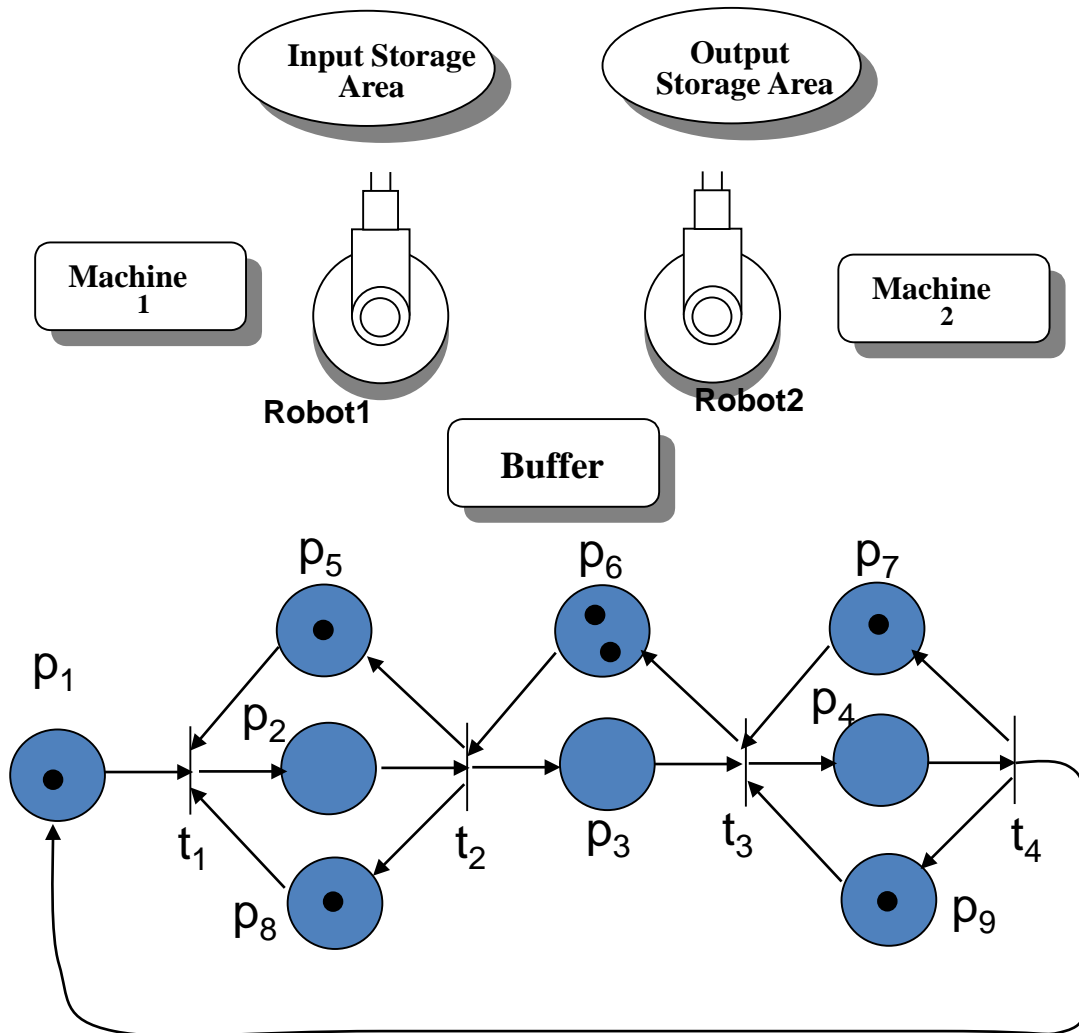
$$F(t_j) = \sum_{m_i \in G} \lambda_j(m_i) \times \pi_i$$

This measures the **system throughput**. Assume that firing each of t_a, t_b, \dots, t_c implies completion of a product, then the entire system's throughput is

$$F(t_a) + F(t_b) + \dots + F(t_c).$$

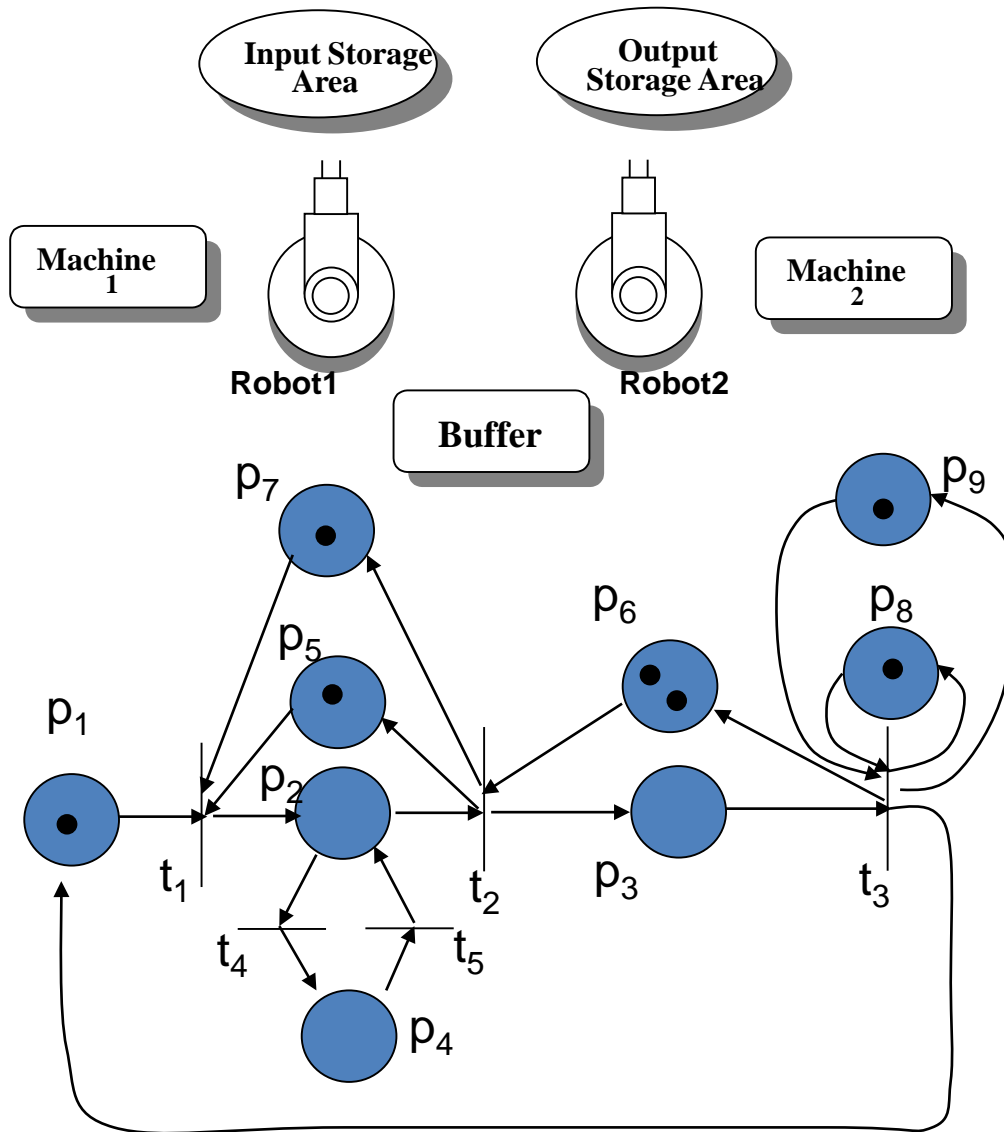
Example: Manufacturing cell

An example we saw earlier in DTPN:



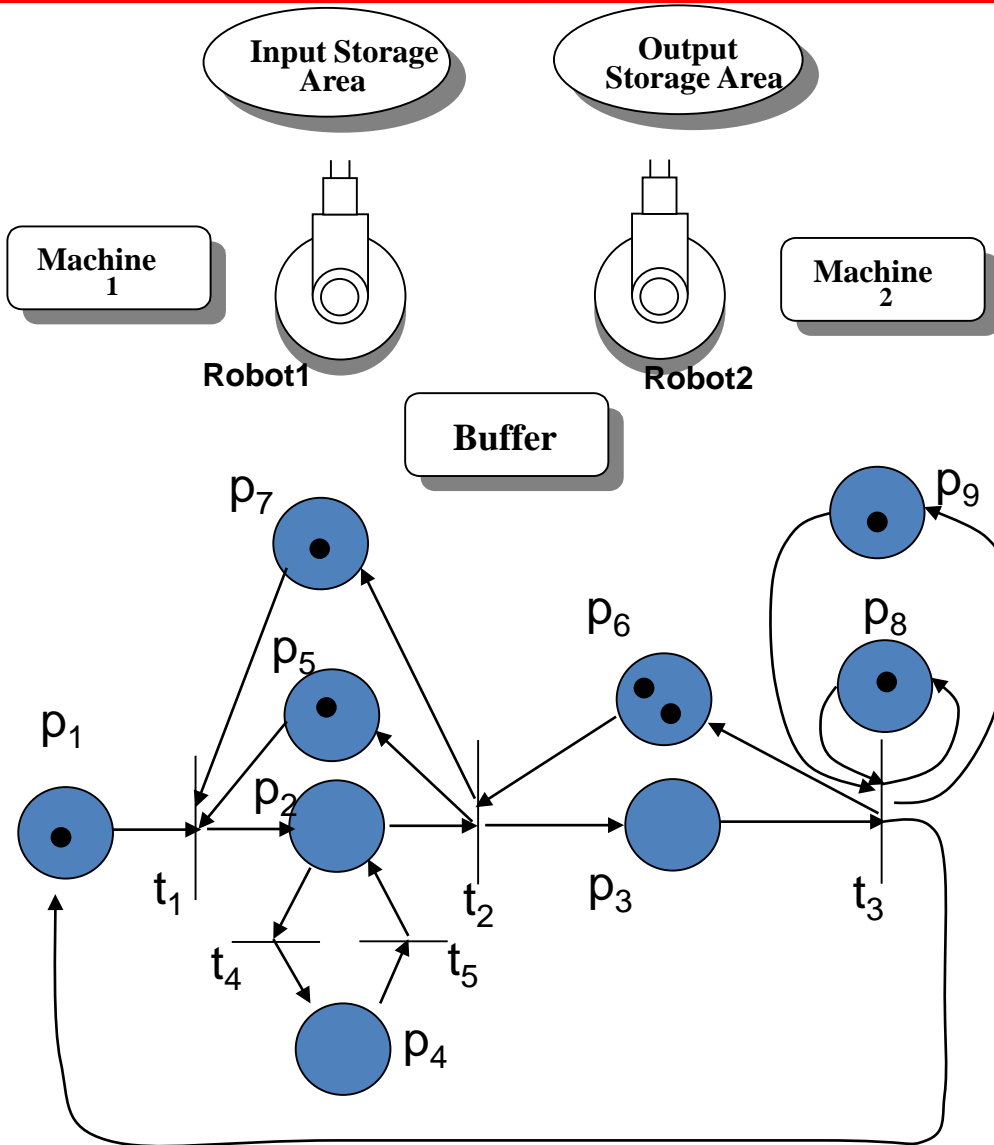
- p_1 : Raw part available
- p_2 : M1 processing a part (**10**)
- p_3 : Intermediate parts ready
- p_4 : M2 processing a part (**16**)
- p_5 : M1 available
- p_6 : buffer slots available
- p_7 : M2 available
- p_8 : R1 available
- p_9 : R2 available
- t_1 : R1 loading (**1**)
- t_2 : R1 unloading (**1**)
- t_3 : R2 loading (**1**)
- t_4 : R2 unloading (**1**)
- (t_4, p_1) : load a raw part to input storage area (**2**)

Example: Manufacturing cell



- M1 may break down. On the average, M1 takes 2 time units to break down and $\frac{1}{4}$ time unit to be repaired. i.e. its average failure and repair rates (1/time unit) are 0.5 and 4.
- M2 is failure-free.
- R1's loading rate is 40. Average rate for M1's processing a part plus R1's unloading is 5 per unit time.
- Average rate for M2's processing plus the R2 loading and unloading is 4 per unit time.

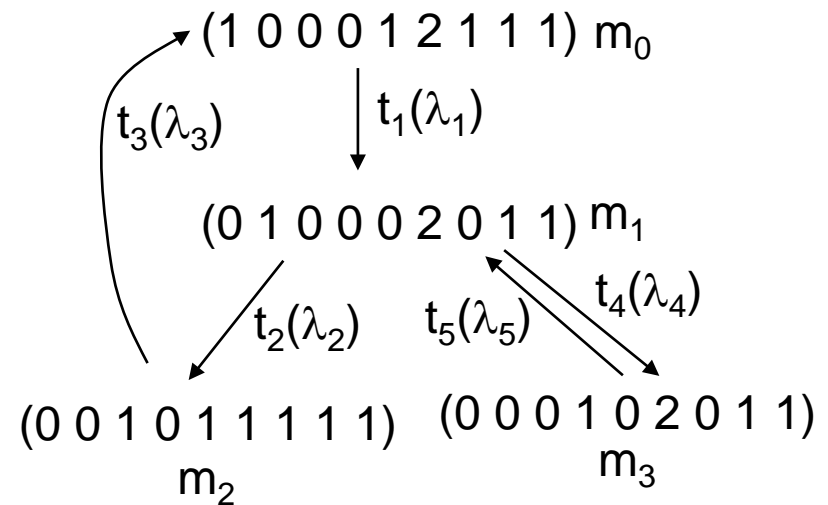
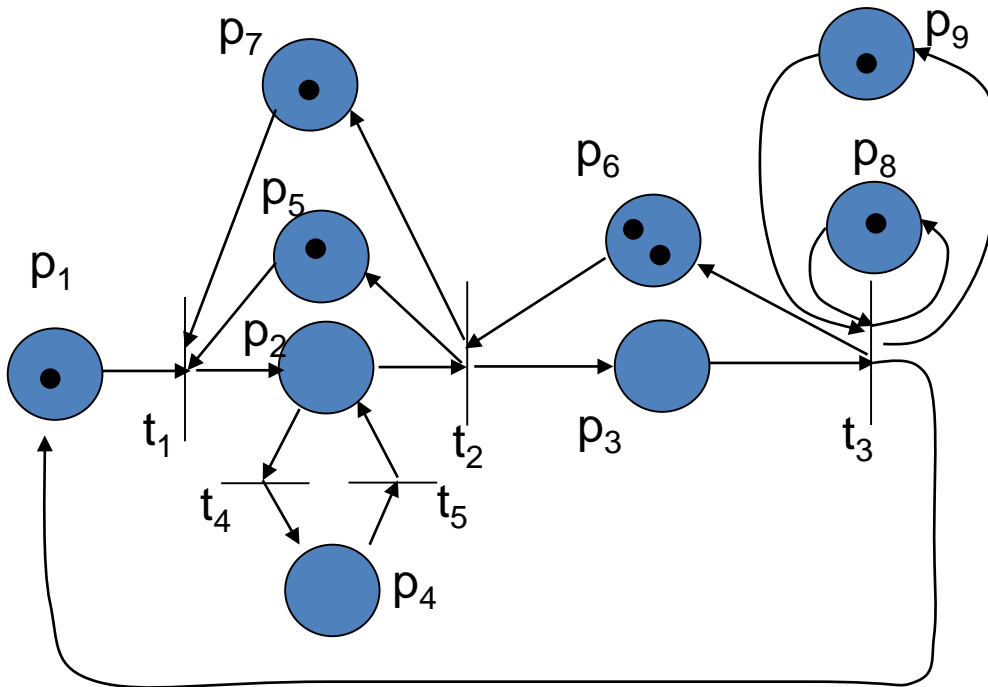
Example: Manufacturing cell



- p_1 : parts are available
- p_2 : M1 is in processing
- p_3 : Part is ready for M2
- p_4 : M1 is in repair
- p_5 : M1 is available
- p_6 : buffer is available
- p_7 : R1 is available
- p_8 : M2 is available
- p_9 : R2 is available
- t_1 : R1 is loading part to M1 (40)
- t_2 : M1 is processing and R2 is unloading (5)
- t_3 : M2 is processing, R2 is loading & unloading (4)
- t_4 : M1 is breaking down (0.5)
- t_5 : repair (4)

Example: Manufacturing cell

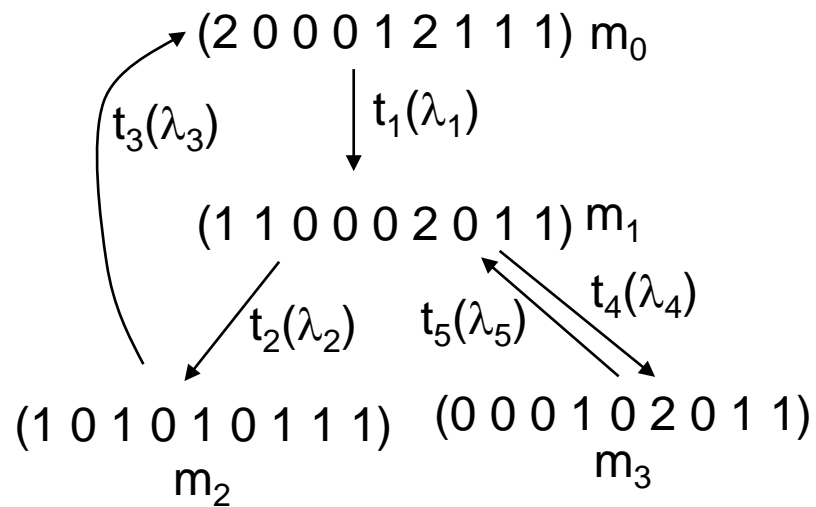
$$\lambda_1=40; \lambda_2=5; \lambda_3=4; \lambda_4=0.5; \lambda_5=4$$



of states: 4

$$A = \begin{bmatrix} -\lambda_1 & \lambda_1 & 0 & 0 \\ 0 & -\lambda_2 - \lambda_4 & \lambda_2 & \lambda_4 \\ \lambda_3 & 0 & -\lambda_3 & 0 \\ 0 & \lambda_5 & 0 & -\lambda_5 \end{bmatrix}$$

What is the average utilization of M1 & system throughput?



Example: Manufacturing cell

Steady-state probabilities can be obtained by solving the following equations:

$$\begin{bmatrix} \pi_0 & \pi_1 & \pi_2 & \pi_3 \end{bmatrix} A = 0 \quad \text{where } A = \begin{bmatrix} -\lambda_1 & \lambda_1 & 0 & 0 \\ 0 & -\lambda_2 - \lambda_4 & \lambda_2 & \lambda_4 \\ \lambda_3 & 0 & -\lambda_3 & 0 \\ 0 & \lambda_5 & 0 & -\lambda_5 \end{bmatrix}$$
$$\pi_0 + \pi_1 + \pi_2 + \pi_3 = 1$$

We obtain:

$$\lambda_1=40; \lambda_2=5; \lambda_3=4; \lambda_4=0.5; \lambda_5=4$$

$$\pi_0 = \lambda_2 \lambda_3 \lambda_5 / \lambda$$

$$\pi_1 = \lambda_1 \lambda_3 \lambda_5 / \lambda$$

$$\pi_2 = \lambda_1 \lambda_2 \lambda_5 / \lambda$$

$$\pi_3 = \lambda_1 \lambda_3 \lambda_4 / \lambda$$

$$\pi_0 = 0.05$$

$$\pi_1 = 0.4$$

$$\pi_2 = 0.5$$

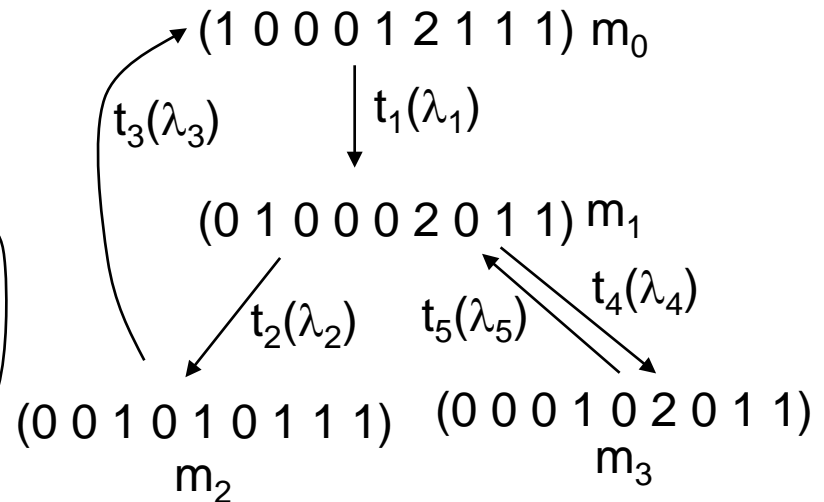
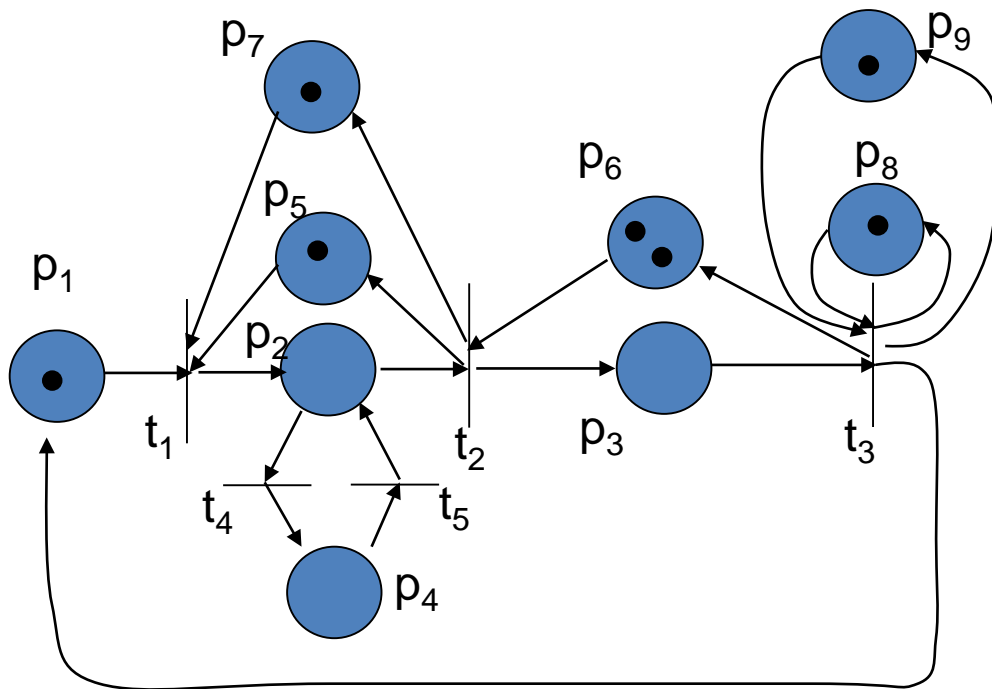
$$\pi_3 = 0.05$$

$$\lambda = \lambda_2 \lambda_3 \lambda_5 + \lambda_1 \lambda_3 \lambda_5 + \lambda_1 \lambda_2 \lambda_5 + \lambda_1 \lambda_3 \lambda_4$$

What is the average utilization of M1 & system throughput?

Example: Manufacturing cell

What is the average utilization of M1?

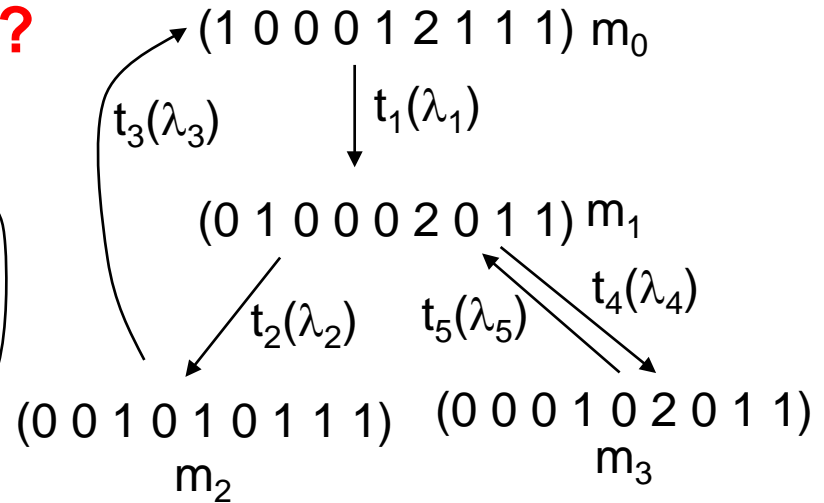
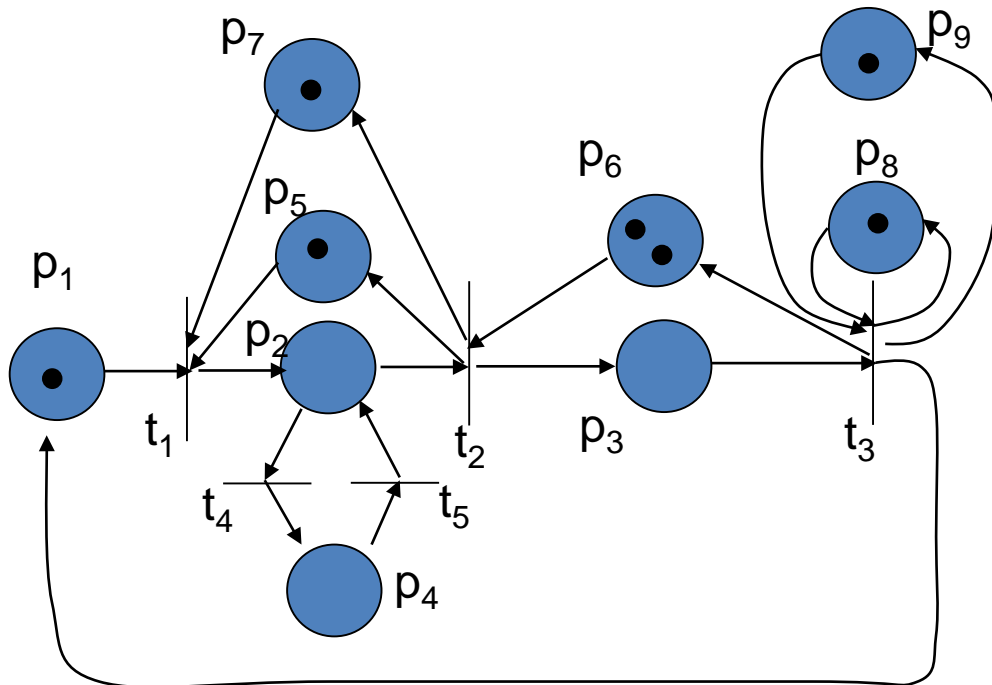


$$\begin{aligned} \pi_0 &= 0.05 \\ \pi_1 &= 0.4 \\ \pi_2 &= 0.5 \\ \pi_3 &= 0.05 \end{aligned}$$

The average utilization of M1 \Leftrightarrow probability that M1 is in process \Leftrightarrow probability of P2 being marked \Leftrightarrow probability of being at marking $m_1 = \pi_1 = 0.4 = 40\%$

Example: Manufacturing cell

What is the system throughput?

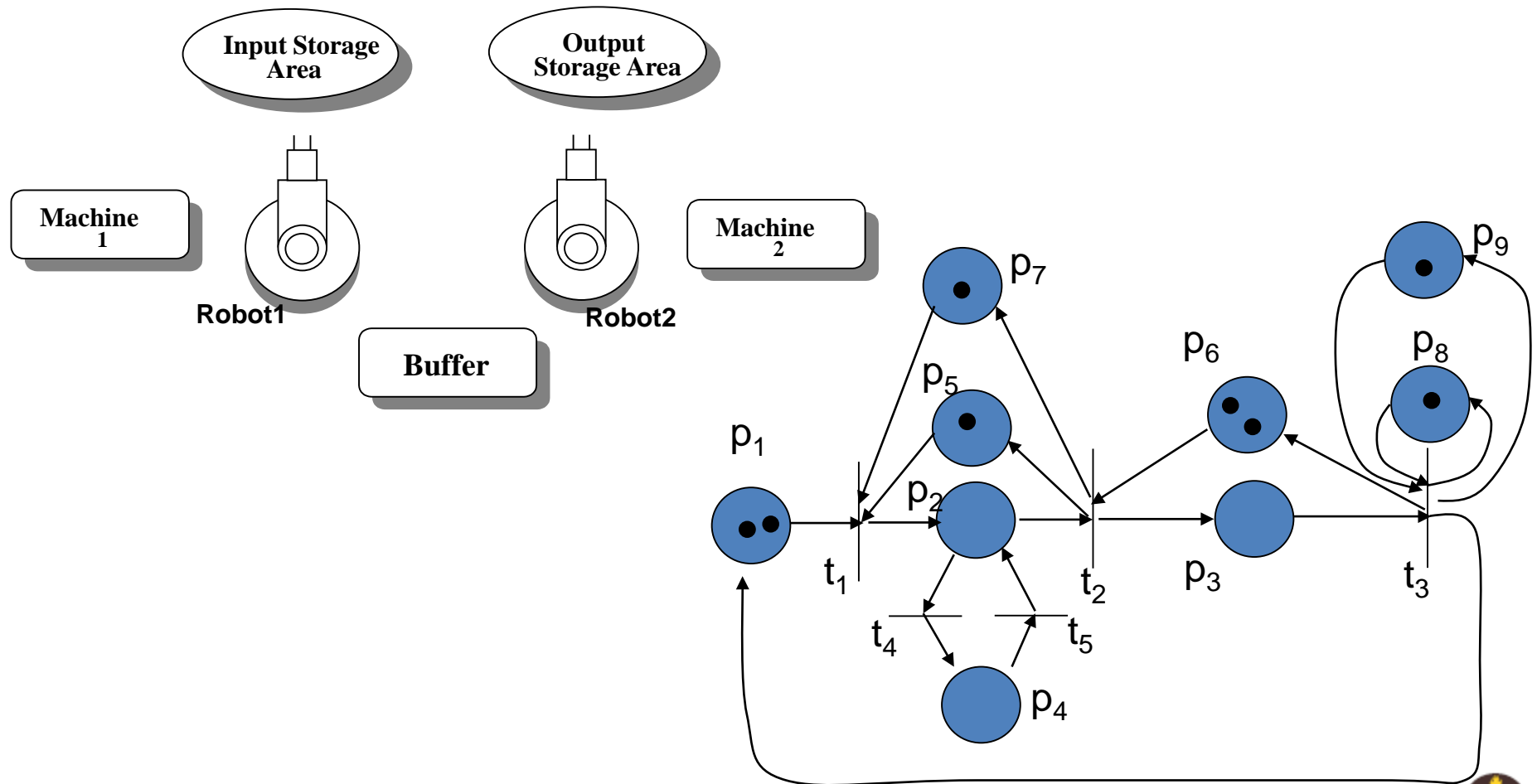


$$\begin{aligned} \pi_0 &= 0.05 \\ \pi_1 &= 0.4 \\ \pi_2 &= 0.5 \\ \pi_3 &= 0.05 \end{aligned}$$

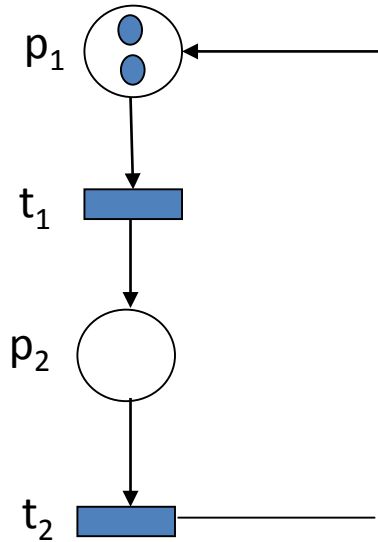
The completion of a product is modeled by firing t_3 . t_3 is enabled at m_2 . System throughput = $\pi_2 * \lambda_3 = 0.5 * 4 = 2$ workpieces/unit time

Exercise: Manufacturing cell

What is the average utilization of M1 and the system throughput if two parts are in the input area?



Marking-dependent transition firing rate



$$m_0 = (2 \ 0)^T$$

$$t_1 \downarrow \uparrow t_2$$

$$m_1 = (1 \ 1)^T$$

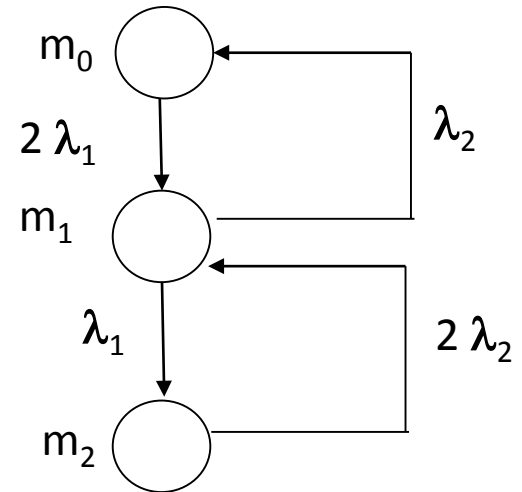
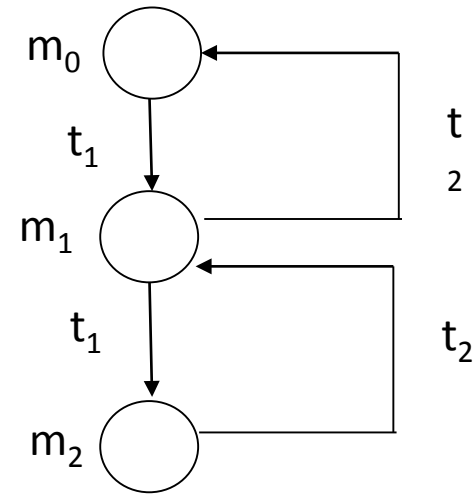
$$t_1 \downarrow \uparrow t_2$$

$$m_2 = (0 \ 2)^T$$

Assume that

t_1 's firing rate is $m(p_1)\lambda_1$

t_2 's firing rate is $m(p_2)\lambda_2$



Marking-dependent transition firing rate

Throughput Performance of the example:

- Solving the balance equations leads to steady state probabilities, π_0 , π_1 and π_2 .
- Then if firing t_2 means producing an item, the throughput is: $\lambda_2\pi_1 + 2\lambda_2\pi_2$
- The average # of tokens in p_1 is $2\pi_0 + \pi_1$
- The average # of tokens in p_2 is $\pi_1 + 2\pi_2$